

Optech Schnorr/Taproot Workshop

September 2019

Welcome!

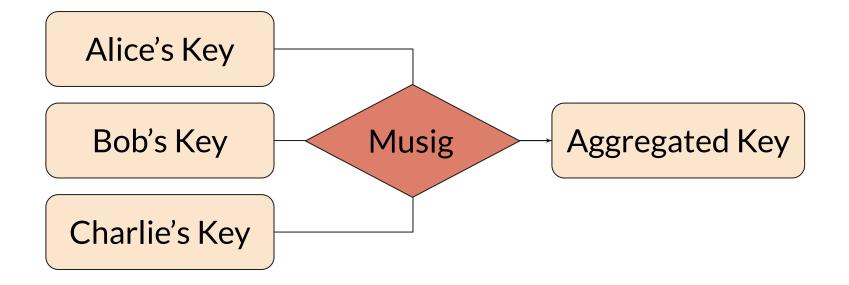
Why Schnorr/Taproot?

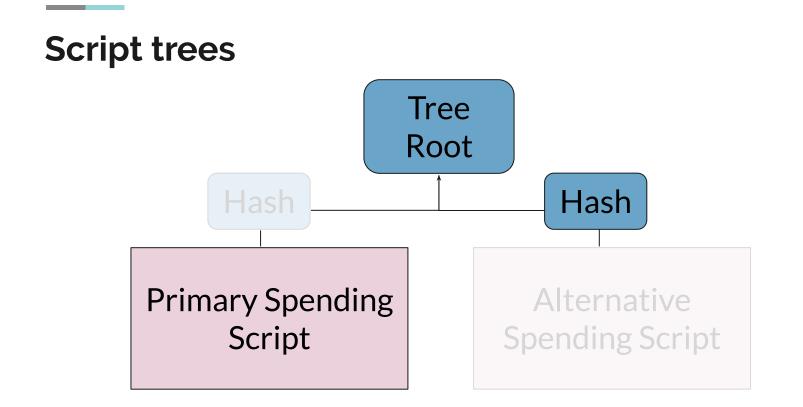
1Scalability• 30-75% savings on multisig
• 2.5x faster block validation2Privacy and Fungibility• All outputs and most spends indistinguishable3Functionality• Very large k of n multisig
• Larger scripts
• Script innovation

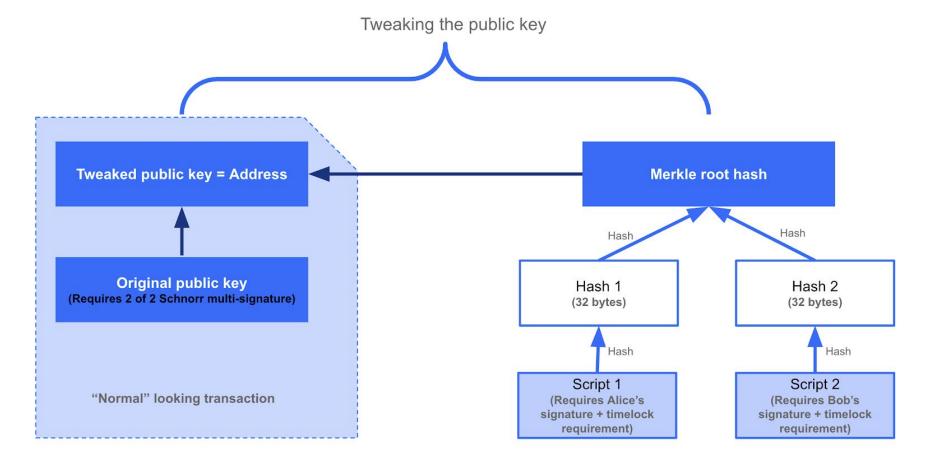
Schnorr signatures

- 1. Better in every way than ECDSA
- 2. 11% smaller than existing signatures
- 3. Compatible with existing private keys
- 4. Same security assumption...with a theoretical proof
- 5. Verification algorithm is linear

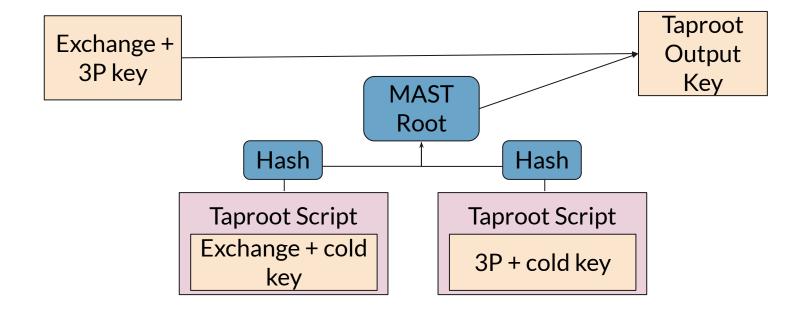
Schnorr enables key and signature aggregation







Exchange 2-of-3 using Musig keytrees



Why Optech?

Bitcoin Optech helps Bitcoin users and businesses integrate scaling technologies.

We provide workshops, documentation, weekly newsletters, original research, case studies and announcements, analysis of Bitcoin software and services, and help facilitate improved relations between businesses and the open source community.

Why this workshop?

- Help share current thinking on schnorr/taproot
- Give engineers a chance to play with the technology
- Involve engineers in the feedback process

WARNING!

The schnorr/taproot proposal is a proposal

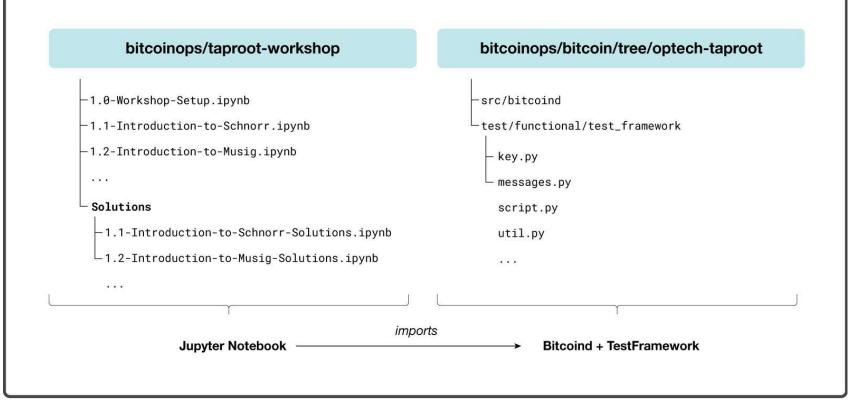
- Details will change
- There is no roadmap
- The workshop code is for educational purposes only!

Chapter 0.1 Toolchain Setup

Did you do your homework?

- Optech Bitcoin Repository: <u>https://github.com/bitcoinops/bitcoin/releases/tag/Taproot_V0.1.4</u>
- Workshop Repository: <u>https://github.com/bitcoinops/taproot-workshop</u>
- Pull latest taproot-workshop
- \$jupyter-notebook
 - 0.1-test-notebook

Optech Schnorr & Taproot Workshop Repositories



Chapter 0.2 Elliptic Curve Math

Scalars (numbers)

- Regular arithmetic but modulo the group order (SECP256K1_ORDER)
- a·b mod n
- Division done using modular inverse (i.e. Fermat's little theorem: $a^p = a$)
- Numbers can go from 0 to (group order 1). eg:
 - \circ (15 + 9) mod 21 = 24 mod 21 = 3
 - (-3) mod 21 = (21-3) = 18

Points on the elliptic curve

- Point = (x, y)
- G is the generator point for our group. (i.e. **P** = dG)
- The curve points form an abelian group:
 - **Closure**: if **A** is a point and **B** is a point than **A** + **B** is a point.
 - Associativity: (A + B) + C = A + (B + C)
 - Identity element: $A + \infty = \infty + A = A$
 - Inverse: For every point A there exist another point B such that A + B = 0
 - Commutativity: A + B = B + A
- Scalar operations:
 - scalar * point: $sG = \{G + G + G + G \dots \text{ s times}\}$
 - point by point division isn't feasible and requires solving **discrete log**

Chapter 1.1 **Schnorr**

Schnorr

Signing:

$$egin{aligned} e &= H(R||P||m) & Sig(s,kG) \ s &= egin{aligned} s &= e \ s &= e \ Sig(s,R) \end{aligned}$$

Verifying:

sG = kG + edG

X only R Points/Public Keys

• Secp256k1:
$$y^2=x^3+7$$

• Solve for y:
$$y=\pm\sqrt{x^3+7}$$

- (-a) mod n = n a
- Even/odd only (odd-even=odd; odd-odd=even)
- Lower/higher half
- Quadratic residue

Chapter 1.2 MuSig

m - message e = H(R||P||m)Naive key aggregation G - generator point d - private key $P_1 = d_1 G, \quad P_2 = d_2 G$ point - scalar * G = (x,y)P - public key (P = dG) $s_1=k_1+ed_1,\ \ s_2=k_2+ed_2$ k - random nonce R - nonce point (**R** = kG) $s_1 + s_2 = (k_1 + k_2) + e(d_1 + d_2)$ s' = k' + ed' $P' = (d_1 + d_2)G$

Glossary

Key cancellation (rogue key) attack $P_1=d_1G, \ P_2=d_2G$ $P_2'=P_2-P_1$ $P'=P_1+(P_2-P_1)$

Glossary m - message e = H(R||P||m)G - generator point d - private key point - scalar * G = (x,y)P - public key (P = dG) k - random nonce R - nonce point ($\mathbf{R} = \mathbf{kG}$)

Musig coefficients $P_1 = d_1 G, \quad P_2 = d_2 G$ $c_i = H(P_1 || P_2 || P_i)$ $d_1' = c_1 d_1, \quad d_2' = c_2 d_2$ $P' = c_1 P_1 + c_2 P_2$

Glossary m - message e = H(R||P||m)G - generator point d - private key point - scalar * G = (x,y)P - public key (P = dG) k - random nonce R - nonce point ($\mathbf{R} = \mathbf{kG}$)

Glossary m - message e = H(R||P||m)Nonce commitments G - generator point d - private key $R_1 = k_1 G$, $R_{2} = k_{2}G$ point - scalar * G = (x,y)P - public key (P = dG) $Com_1 = H(R_1), \ Com_2 = H(R_2)$ k - random nonce R - nonce point ($\mathbf{R} = \mathbf{kG}$) $R' = R_1 + R_2$ e = H(R'||P'||m)

Chapter 2.1 - 2.4 **Taproot**

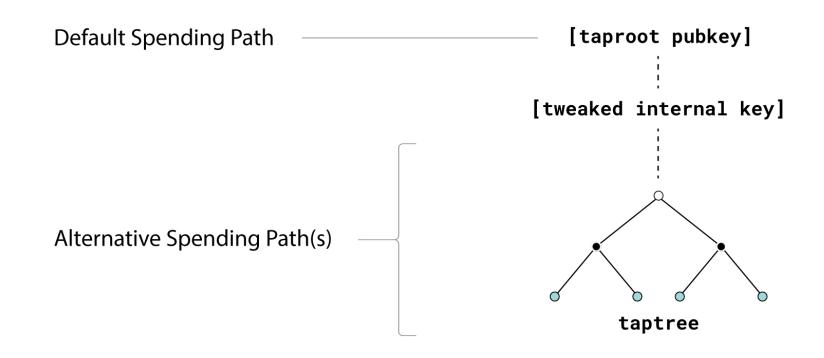
Default & Alternative Spending Paths

- Default Spending Path
 - Single or multi-party public keys (indistinguishable)
- Alternative Spending Path(s)
 - Single or multiple "hidden" alternative scripts.
 - Only the script of the spent path is revealed when spent.

Taproot: Multi-party contract

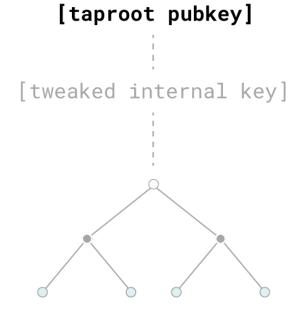
• Default Spending Path

- Aggregated pubkey/signature.
- Default spending path hides multi-party contract.
- Alternative Spending Path(s)
 - In aggregate, enforce the multi-party contract.
 - script_0 OR script_1 OR script_2 ...



Chapter 2.1 Segwit Version 1

Segwit version 1



taptree

Segwit version 1

- Output script:
 - Script: [01] [33B public key]
 - Has recently been reduced to 32B public key in bip-schnorr.
 - This workshop has been built with the previous 33B public key format.
- Satisfying Witness:

Ο

Key path:

[64B BIP-schnorr signature]

• Script path: [initial stack] [tapscript] [controlblock]

P2PK vs P2PKH

- P2PK vs P2PKH:
 - V1 Script:
 - V1 Witness:
 - V0 Script:
 - V0 Witness:
- V1 program witness:
- Disadvantages of P2PKH:
 - Cost:

[01] [33B public key]
[64/65B signature]
[00] [20B pubkey hash]
[DER signature(ecdsa)] [public key]
single key, MuSig, ...

pubkey + pubkey hash

Taproot Sighash Flags

- Taproot retains legacy sighash flag semantics
 - ALL, NONE, SINGLE, ANY
 - New implied ALL sighash flag (0x00)

Taproot: Schnorr signature encoding

1B



- x(R), s, sighashflag
 - x(R): 32B
 s: 32B
 - SIGHASH flag:

(All, None, Single, Any) 0x01, 0x02, 0x03, 0x8...

v1: schnorr signature hash

- Control
 - Always epoch(0) | sighash
- Transaction
 - Always
 - $\circ \quad \text{If !any} \quad$
 - If !none or !single

• Input

- \circ Always
- $\circ \quad \text{If any} \quad$
- If !any
- Output(s)
 - \circ If single

sha256(CTxOut)

input index

version | locktime

spend_type | scriptPubKey

outpoint | input amount | sequence

outputs

prevout(s) | input amount(s) | sequence(s)

v1: schnorr signature hash

•	Control		
	 Always 	epoch(0) sighash	
•	Transaction		Deveeble Midstate
	 Always 	version locktime	Reusable Midstate
	 If !any 	prevout(s) input amount(s) sequence(s)	
	• If !none or !sing	le outputs	
•	Input		
	 Always 	spend_type scriptPubKey	
	• If any	outpoint input amount sequence	
	 If !any 	input index	
•	Output(s)		
	• If single	sha256(CTxOut)	

Chapter 2.2 **Taptweak**

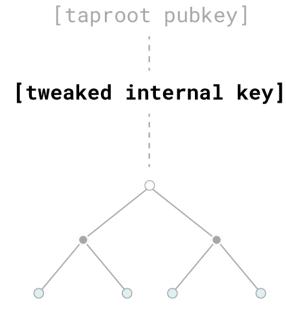
Taptweak

- Any data can be committed to a public key tweak.
 - Public key remains spendable.
 - Owner of private key can spend with knowledge of tweak.
 - Signing with the tweaked public key does not reveal tweak.
 - The owner of the private key can later reveal the commitment without revealing the private key.

TapTweak

- A TapTweak is a tweak to an internal public key
 - Default spending path: Tweak is not revealed.
 - Alternative spending paths:
 - Tweak & script branch are revealed.
 - Script branch is executed during validation.

Taptweak



taptree

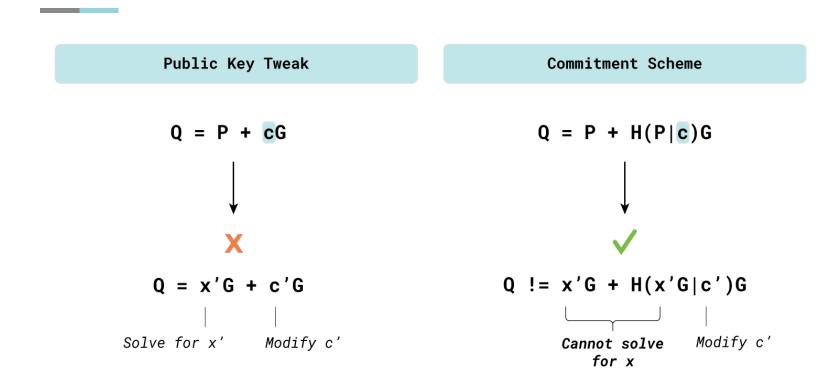
Committing data to a pubkey tweak

• v1 witness program : 33B pubkey **Q**

Q = **P** + **H**(**P**|**c**)**G** where **P** is the *internal key* and **c** is the *commitment*.

• Spending witness: 64B signature (x(R), s)

The private key is tweaked with **H(P|c)** before signing

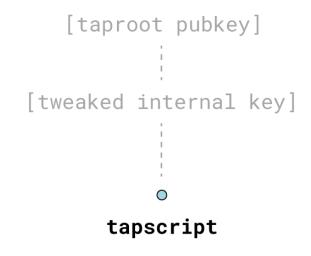


Chapter 2.3 **Tapscript**

Tapscript

- TapScript is upgraded Bitcoin script.
 - Optimized for Schnorr.
 - Allows for future TapScript versions.
 - TapScripts are committed to TapTweaks.

Tapscript



Tapscript vs. Bitcoin script

- Signature opcodes: Perform verification of bip-schnorr signatures
- Multisig opcodes: Removed
- Checksigadd opcodes: **Replace multisig opcodes. Enable signature batch verification.**
- Versioning:
 - TapLeaf version: **0xc0**
 - Upgradable opcodes: **80, 98, 126-129, 131-134, 137-138, 141-142, 149-153, 187-254**
 - Difference to NOP: Immediate success and termination of script execution.

Multisig with Checksigadd

- Output Script
- Initial Stack
- pk0 sig0
- checksig sig1
- pk1 sig2
- \circ checksigadd
- o pk2
- checksigadd
- o **3**
- equal

Multisig with Checksigadd

- Output Script
 Initial Stack
 - pk1 **1**
 - \circ checksigadd \circ sig1
 - pk2 sig2
 - checksigadd
 - o **3**
 - equal

Multisig with Checksigadd

• Output Script

0

• Initial Stack

3

0

• equal

3

Tapscript Descriptors (I/II)

- Pay-to-pubkey:
 - ts(pk(key))
 - ts(pkhash(key, digest))
 - ts(pkolder(key, delay))
 - ts(pkhasholder(key, digest, delay))

Satisfying Witness:

[signature]

[preimage] [signature]

[signature] (nSequence > delay)

[preimage][signature] (nSequence > delay)

Tapscript Descriptors (II/II)

- Pay-to-pubkey:
 - ts(csa(k, keys..))
 - ts(csahash(k, keys, digest))
 - ts(csaolder(k, keys, delay))
 - ts(csahasholder(k, keys, digest, delay))

Satisfying Witness: [k signatures] [hash] [k signatures] [k signatures] (nSequence > delay) [hash] [k signatures] (nSequence > delay)

Committing a single Tapscript to a Taptweak

- Taptweak t
 - **Q = P + tG**
 - t = TaggedHash("TapTweak", P, tapleaf)
 - TapLeaf = TaggedHash("TapLeaf", ver, size, script)
- TaggedHash
 - TaggedHash(data) = sha256(sha256("Tag") + sha256("Tag") + data)
 - Collision resistance
 - 64B re-usable midstate

Taproot: Spending a single Tapscript

Spending Witness:

- [Satisfying witness elements for Tapscript]
- [Tapscript]
- [Internal Key]

Unspendable script path (WIP)

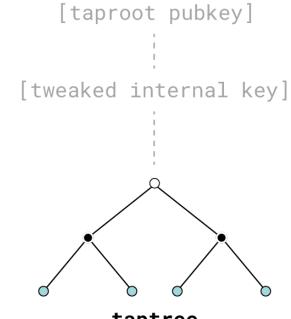
- Problem: Hidden script path t'
 - Q = P1 + P2 = P1 + P2' + H(P1+P2'|t')
- Solution: Default unspendable script path t
 - $\circ \qquad \mathsf{Q} = \mathsf{P1} + \mathsf{P2} + \mathsf{H}(\mathsf{P1} + \mathsf{P2}|\mathbf{t})\mathsf{G}$
 - Not possible:
 - Hidden t': P2 = P2' + H(Pagg|t')
 - Default t: Q = P1 + P2 + H(P1+P2|t)

Chapter 2.4 **Taptree**

Taptree

- A Taptree commits multiple Tapscripts to a Taptweak
 - Binary merkle tree commitment structure.
 - A TapTree does not have to be balanced.
 - Allows for Tapscript specific spending cost optimizations.

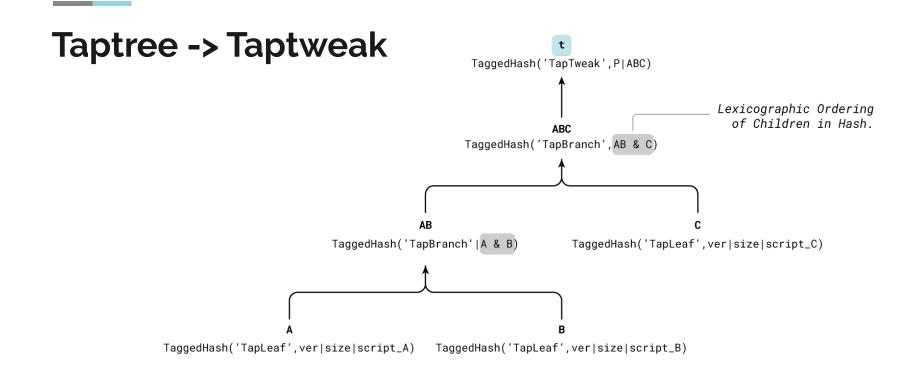
Taptree



taptree

Committing Tapscripts to a Taptweak

- TapTweak t
 - **Q = P + tG**
 - t = TaggedHash("TapTweak", P, Tapbranch)
 - Tapbranch is the root node of the TapTree



Committing Tapscripts to a Taptweak

- TapTweak t
 - **Q = P + tG**
 - o t = TaggedHash("TapTweak", P, Root)
 - Root is root node of TapTree
- TapTree
 - Binary tree
 - Siblings ordered lexicographically
 - Internal nodes are tagged "TapBranch"
 - Leaf nodes are tagged "TapLeaf"
 - TapScripts are committed to leaf nodes

Protects against preimage attacks.

Taproot Descriptors

- Taproot Descriptor:
 - P = Internal Pubkey
 - Tweak is implied from taptree descriptor
- Taptree Descriptor:
 - TapBranch represented by
 - TapBranch are composable

[tapscript0, [tapscript1, tapscript 2]] [child_node0, child_node1] [tapscript0, [tapscript1, tapscript2]]

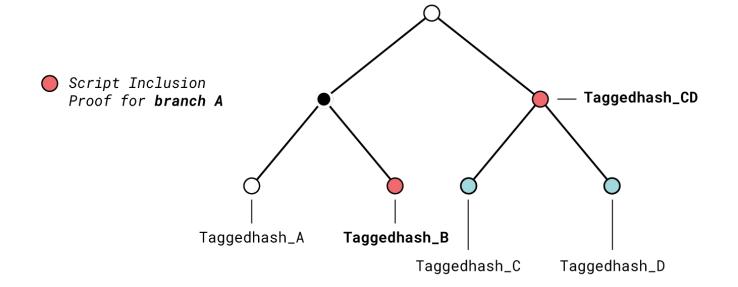
tp(P, [taptree descriptor])

Taproot: spending a script path

- Taproot descriptor:
- Satisfying witness for script 1:
- [scrip

— Inclusion proof for script A





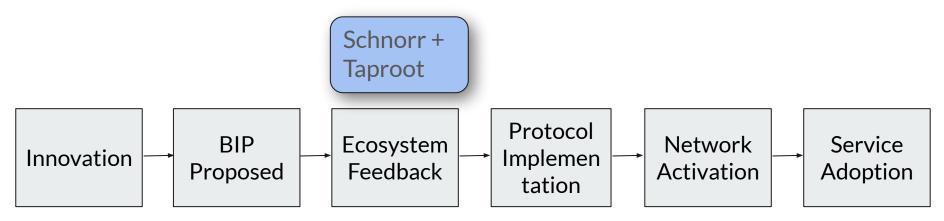
Chapter 3.1 Case Study

Discussion

Where to find out more

- Draft BIPs: <u>https://github.com/sipa/bips/tree/bip-schnorr</u>
- Reference implementation: <u>https://github.com/sipa/bitcoin/tree/taproot</u>
- Mailing list: <u>https://lists.linuxfoundation.org/pipermail/bitcoin-dev/</u>

Bitcoin Consensus Upgrade Lifecycle



Mailing list

[bitcoin-dev] Taproot proposal

Pieter Wuille pieter.wuille at gmail.com

Mon May 6 17:57:57 UTC 2019

- Previous message: [bitcoin-dev] Bitcoin Knots 0.18.0.knots20190502 released
- Next message: [bitcoin-dev] Taproot proposal
- Messages sorted by: [date] [thread] [subject] [author]

Hello everyone,

Here are two BIP drafts that specify a proposal for a Taproot softfork. A number of ideas are included:

 \ast Taproot to make all outputs and cooperative spends indistinguishable from eachother.

- * Merkle branches to hide the unexecuted branches in scripts.
- * Schnorr signatures enable wallet software to use key aggregation/thresholds within one input.

* Improvements to the signature hashing algorithm (including signing all input amounts).

* Replacing OP_CHECKMULTISIG(VERIFY) with OP_CHECKSIGADD, to support batch validation.

* Tagged hashing for domain separation (avoiding issues like CVE-2012-2459 in Merkle trees).

 \ast Extensibility through leaf versions, OP_SUCCESS opcodes, and upgradable pubkey types.

- [bitcoin-dev] Taproot proposal Pieter Wuille
 - [bitcoin-dev] Taproot proposal Luke Dashjr
 - [bitcoin-dev] Taproot proposal Sjors Provoost
 - [bitcoin-dev] Taproot proposal ZmnSCPxj
 - [bitcoin-dev] Taproot proposal ZmnSCPxj
 - [bitcoin-dev] Taproot proposal Pieter Wuille
 - [bitcoin-dev] Taproot proposal ZmnSCPxj
 - [bitcoin-dev] Taproot proposal ZmnSCPxj
 - [bitcoin-dev] Taproot proposal Johnson Lau
 - [bitcoin-dev] Taproot proposal ZmnSCPxj
 - [bitcoin-dev] Taproot proposal Anthony Towns
 - [bitcoin-dev] Taproot proposal Luke Dashjr
 - [bitcoin-dev] Taproot proposal Russell O'Connor
 - [bitcoin-dev] Taproot proposal Pieter Wuille
 - [bitcoin-dev] Taproot proposal Russell O'Connor
 - [bitcoin-dev] Taproot proposal John Newbery



Why this workshop?

- Help share current thinking on schnorr/taproot
- Give engineers a chance to play with the technology
- Involve engineers in the feedback process

Contributions welcome!

https://github.com/bitcoinops/taproot-workshop